

How to peel fully convex digital sets ?

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University of Groningen

How to peel fully convex digital sets ?

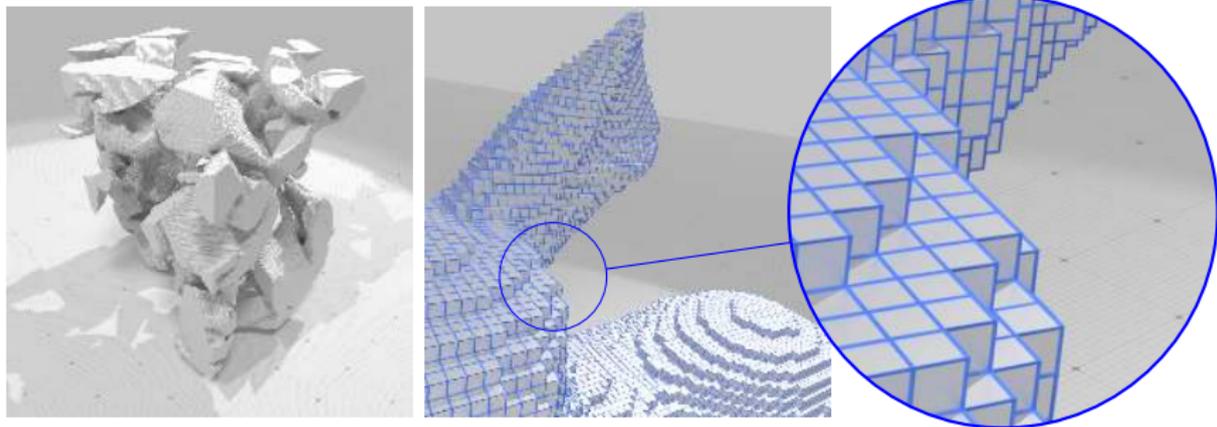
Context and objectives

Sufficient conditions for peelability

Characterization of peelability

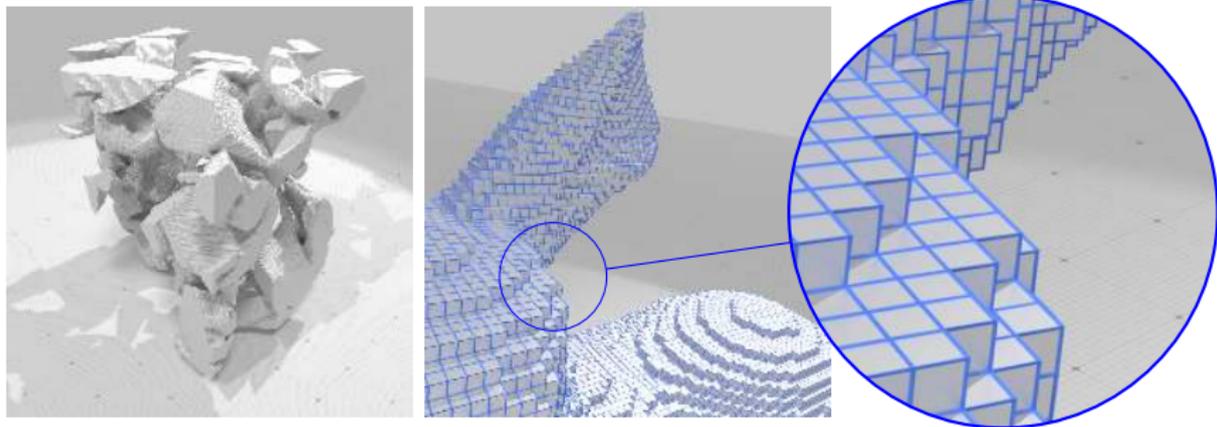
Conclusion and future works

Why digital convexity ?



- ▶ no (infinitesimal) differential geometry for digital shapes
- ▶ convexity: a fundamental tool to analyze the geometry of shapes
- ⇒ convex/concave/flat/saddle regions, piecewise linear geometry

Why digital convexity ?

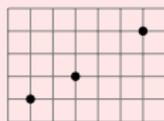


- ▶ no (infinitesimal) differential geometry for digital shapes
 - ▶ convexity: a fundamental tool to analyze the geometry of shapes
- ⇒ convex/concave/flat/saddle regions, piecewise linear geometry
- ▶ convexity = foundation of convex analysis, linear programming
 - ▶ digital convexity = foundation of digital convex analysis, integer linear programming ?

Full convexity versus classical digital convexity ?

Classical digital convexity does not imply **connectedness** !

► $X \subset \mathbb{Z}^d$ is 0-convex iff $\text{Cvxh}(X) \cap \mathbb{Z}^d = X$

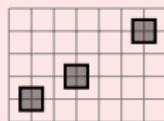


X

=



$\text{Cvxh}(X) \cap \mathbb{Z}^d$



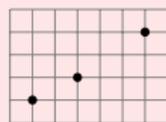
pixel view

0-convex $\not\Rightarrow$ **connected**

Full convexity versus classical digital convexity ?

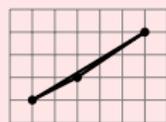
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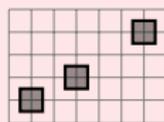


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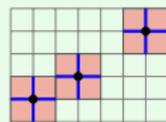


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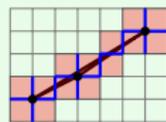
Full convexity implies **digital** and **simple connectedness**

- $X \subset \mathbb{Z}^d$ is *fully convex* iff $\text{Cvxh}(X) \cap \bar{\mathcal{C}}^d = X \cap \bar{\mathcal{C}}^d$, where $\bar{\mathcal{C}}^d$ is the unit cellular grid decomposition of \mathbb{R}^n [L. 2021].

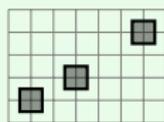


$X \cap \bar{\mathcal{C}}^d$

\neq



$\text{Cvxh}(X) \cap \bar{\mathcal{C}}^d$



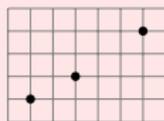
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0-convex but not *fully convex* \Rightarrow **not connected**

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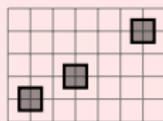


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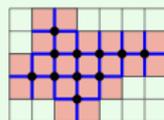


pixel view

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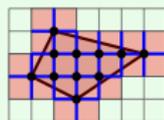
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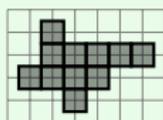


$X \cap \bar{\mathcal{C}}^d$

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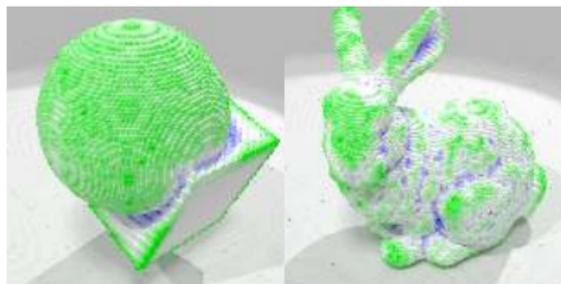
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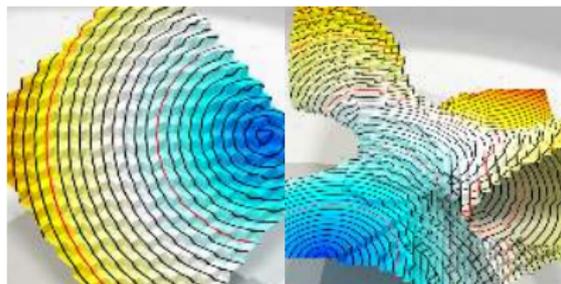
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fully convex \Rightarrow **digitally and simply connected**

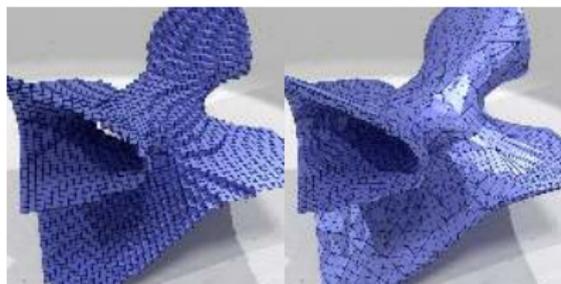
Applications of full convexity to digital shape analysis



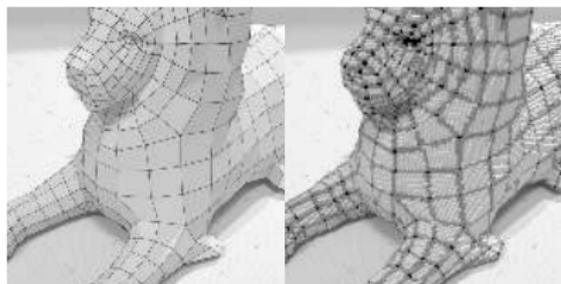
exact local shape analysis
(convex, concave, planar (white))



geodesics
(Euclidean distance in digital planes)

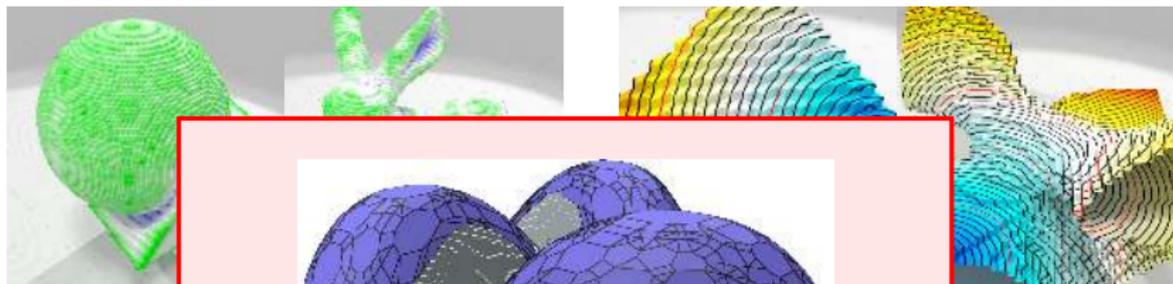


polyhedrization
(Hausdorff 1-close and reversible)



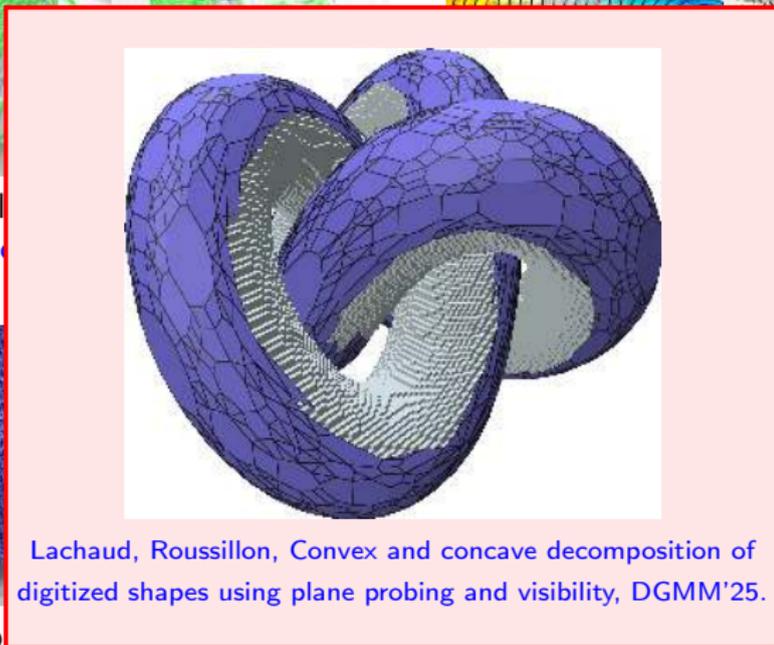
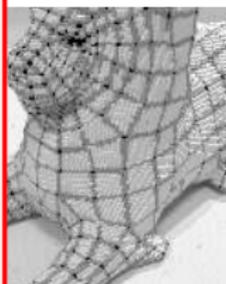
digital polyhedron
(cells are fully convex)

Applications of full convexity to digital shape analysis



exact l
(convex, c

digital planes)

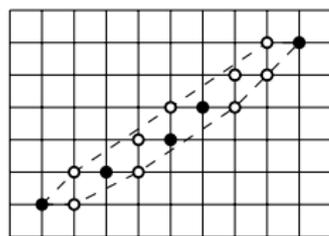


Lachaud, Roussillon, Convex and concave decomposition of digitized shapes using plane probing and visibility, DGMM'25.

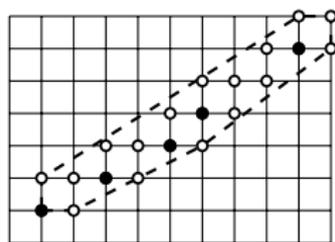
p
(Hausdorff 1-close and reversible)

iron
(cells are fully convex)

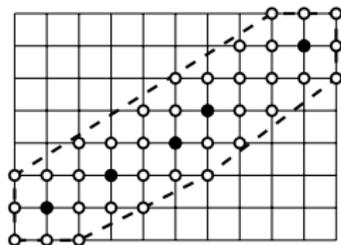
But is there a good fully convex hull operator ?



$FC^*(X)$



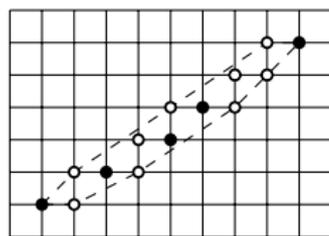
$(Cvxh(X) \oplus H^+) \cap \mathbb{Z}^d$



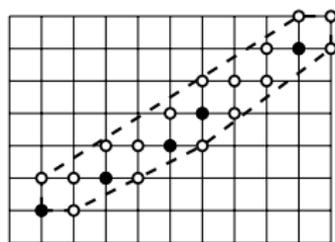
$Extr(Star(Cvxh(X)))$

operator	$FC^*(X)$	$(Cvxh(X) \oplus H^+) \cap \mathbb{Z}^d$	$Extr(Star(Cvxh(X)))$
fully conv.	yes	yes	yes
idempotence	yes	no	no
symmetry	yes	no	yes
$\#(Out)/\#(In)$	low	medium	high
efficiency	iterative	yes	yes
Increasing	not always	yes	yes

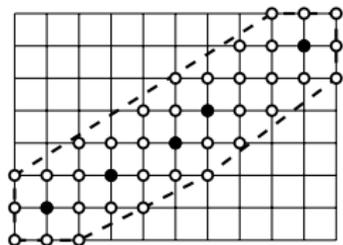
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$FC^*(X)$



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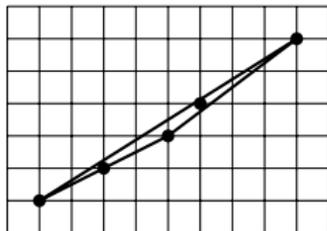
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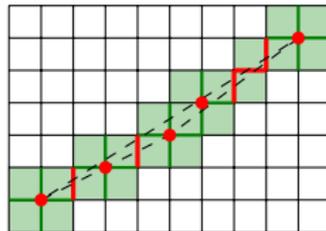
Best compromise: fully convex envelope operator $FC^*(X)$...
but iterative and sometimes not an increasing operator.

Fully convex envelope $FC^*(X)$

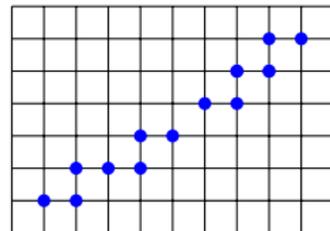
- ▶ Iterative method for computing a fully convex envelope [Feschet, L. '23]
- ▶ Let $FC(X) := \text{Extr}(\text{Skeleton}(\text{Star}(\text{Cvxh}(X))))$
- ▶ Iterative composition $FC^n(X) := \underbrace{FC \circ \dots \circ FC}_n(X)$
- ▶ *Fully convex envelope* of X is $FC^*(X) := \lim_{n \rightarrow \infty} FC^n(X)$.



input X , $Y := \text{Cvxh}(X)$



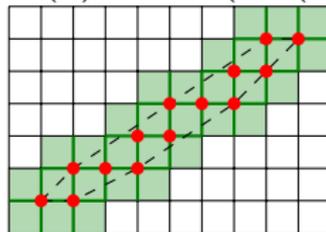
$\text{Star}(Y)$, $\text{Skeleton}(\text{Star}(Y))$



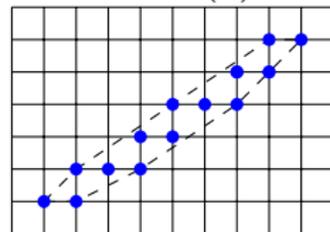
$X' = FC(X)$



input X' , $Y' := \text{Cvxh}(X')$

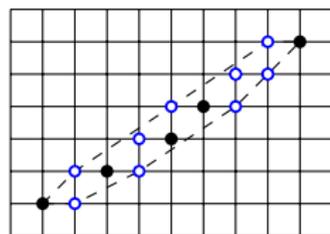


$\text{Star}(Y')$, $\text{Skeleton}(\text{Star}(Y'))$

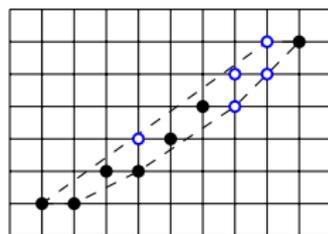


$X'' = FC(X') = FC^2(X)$

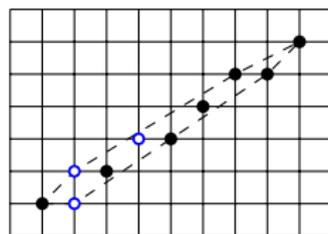
Cons of fully convex envelope $FC^*(X)$?



$X_1, FC^*(X_1)$



$X_2, FC^*(X_2)$



$X_3, FC^*(X_3)$

- ▶ $X_1 \subset X_2$ and $FC^*(X_2) \subset FC^*(X_1)$
- ▶ $X_1 \subset X_3$ and $FC^*(X_3) \subset FC^*(X_1)$
- ▶ $FC^*(\cdot)$ is **not increasing** in general for *thin* sets !
- ▶ $FC^*(\cdot)$ is **increasing** for *thick* sets [Feschet, L. 2023], like $(X \oplus [-1, 1]^d) \cap \mathbb{Z}^d$ for any digital set X .
- ▶ also **hard** to bound the number of iterations of $FC^*(\cdot)$

Peeling fully convex sets

Objective of the paper

Better understanding of the **inclusion hierarchy** of *fully convex sets* in arbitrary dimension d .

Method

- ▶ **Peeling a fully convex set X :**
find the $x \in X$ such that $X \setminus \{x\}$ is still fully convex.

Similar works

- ▶ contractions and dilations of digital convex sets
[Tarsissi *et al.*, '19, '22, '23]
 - ▶ arithmetic approach focused on 2D convexity
- ▶ reminiscent of **simple points** detection
[Bertrand, Couprie, Kong, and many others since the '90]
 - ▶ one of the most successful results of digital topology
 - ▶ peelable points preserve **convexity and topology**

How to peel fully convex digital sets ?

Context and objectives

Sufficient conditions for peelability

Characterization of peelability

Conclusion and future works

“Peeling” conditions for digitally convex sets

Lemma 1 (vertices are peelable points of 0-convex sets)

Let $X \subset \mathbb{Z}^d$ be a 0-convex set, i.e. $\text{Cvxh}(X) \cap \mathbb{Z}^d = X$. Then, for any subset of vertices Z of $\text{Cvxh}(X)$, $X \setminus Z$ is 0-convex.

Lemma 2 (a peelable point is a vertex of $\text{Cvxh}(X)$)

Let $X \subset \mathbb{Z}^d$ be a 0-convex set. Let $z \in X$. Then $X \setminus \{z\}$ is 0-convex if and only if z is a vertex of $\text{Cvxh}(X)$.

Conclusion

Peelable points of a fully convex set are among its vertices.

[Reminder] Full convexity is equivalent to P -convexity

Let π_j be the orthogonal projector associated to the j -th axis.

Projection convexity (P -convexity)

$X \subset \mathbb{Z}^d$ is P -convex (in \mathbb{Z}^d) iff:

- (i) X is digitally 0-convex,
- (ii) when $d > 1$, $\forall j, 1 \leq j \leq d$, $\pi_j(X)$ is P -convex (in \mathbb{Z}^{d-1}).

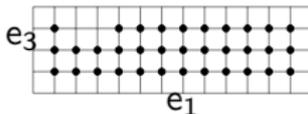
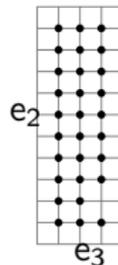
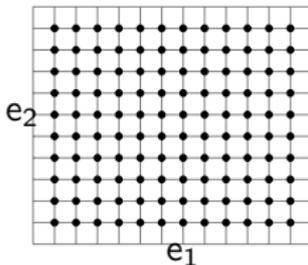
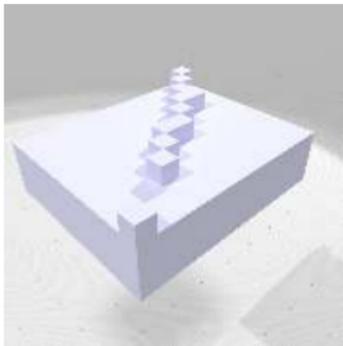
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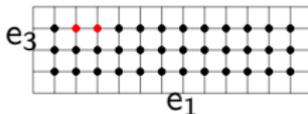
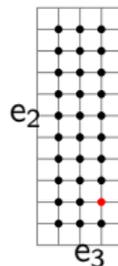
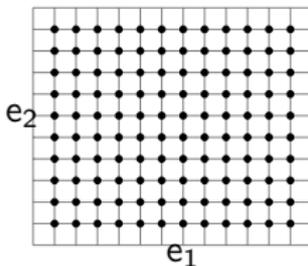
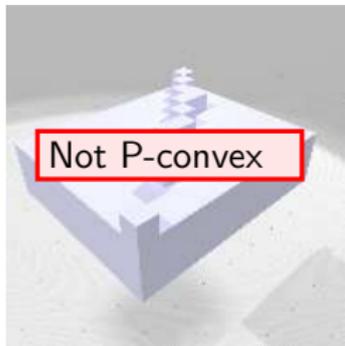
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Not 0-convex

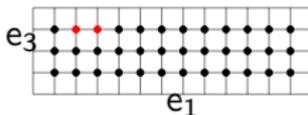
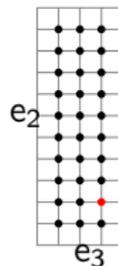
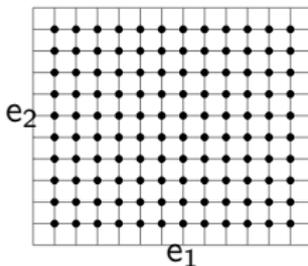
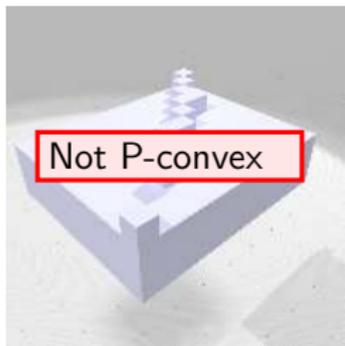
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Not 0-convex

Theorem [Feschet, L. '24]

P -convexity is equivalent to full convexity.

Sufficient conditions for peelability I

Theorem (sufficient conditions)

Let X , finite, non-empty, fully convex set of \mathbb{Z}^d . Let Z be a subset of the vertices of $\text{Cvxh}(X)$.

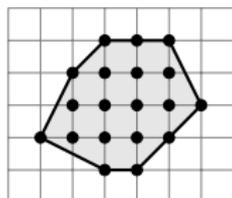
Assume that, for every direction $i \in I_d$, for every $z \in Z$, we have $\{z - e_i, z + e_i\} \cap (X \setminus Z) \neq \emptyset$. Then $X \setminus Z$ is fully convex.

Sufficient conditions for peelability I

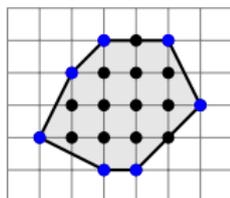
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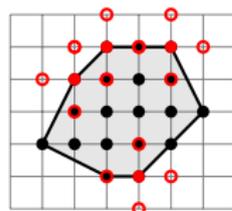
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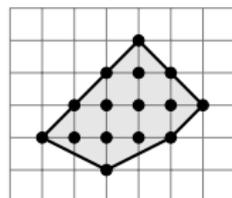
X and $\text{Cvxh}(X)$



$\text{Extr}(\text{Cvxh}(X))$



Z and $Z \pm e_i$ as \circ



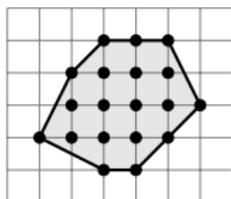
$X \setminus Z$

Sufficient conditions for peelability I

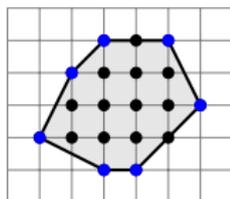
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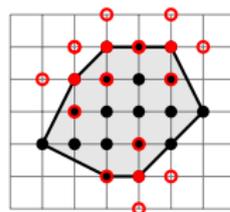
Assume that, for every direction $i \in I_d$, for every $z \in Z$, we have $\{z - e_i, z + e_i\} \cap (X \setminus Z) \neq \emptyset$. Then $X \setminus Z$ is fully convex.



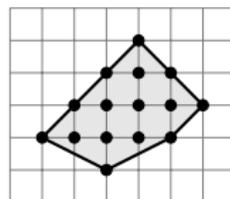
X and $\text{Cvxh}(X)$



$\text{Extr}(\text{Cvxh}(X))$



Z and $Z \pm e_i$ as \circ



$X \setminus Z$

Proof.

1. $X \setminus Z$ is 0-convex because $Z \subset \text{Extr}(\text{Cvxh}(X))$ (Lemma 1).
2. Since $\{z - e_i, z + e_i\} \cap (X \setminus Z) \neq \emptyset$, then $\pi_i(X \setminus Z) = \pi_i(X)$.
3. So $\pi_i(X \setminus Z)$ is P -convex and use equivalence with full convexity.



Sufficient conditions for peelability II

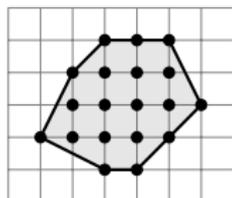
Proposition (side removal)

Let X be any fully convex set of \mathbb{Z}^d . Let $H = \{x \in \mathbb{R}^d, x \cdot e_i \leq \alpha\}$ be any half-space of axis-aligned normal vector e_i and intercept α . Then $X \setminus H$ is fully convex.

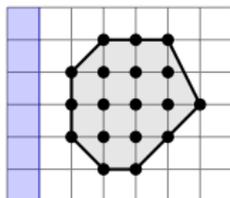
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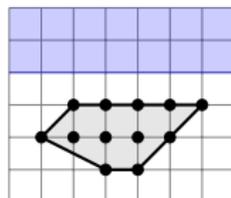
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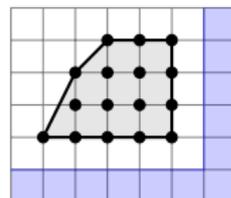
X and $\text{Cvxh}(X)$



$X \setminus H$



$X \setminus H$

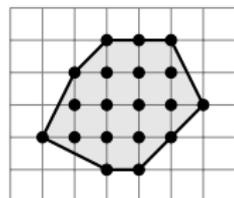


$X \setminus (H \cup H')$

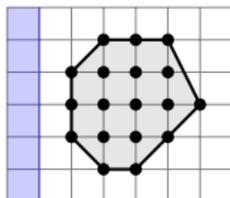
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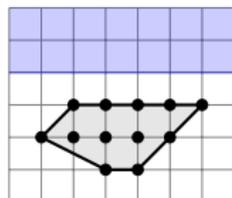
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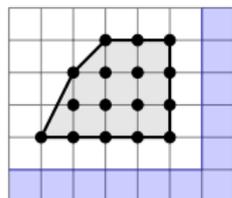
X and $\text{Cvxh}(X)$



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$X \setminus H$



$X \setminus (H \cup H')$

Proof.

Half-space $H^c = \{x \in \mathbb{R}^d, x \cdot e_i \geq \alpha + 1\}$ is *stable*, i.e. it is convex and for any cell σ of \mathbb{C}^d , $H^c \cap \sigma \neq \emptyset \Rightarrow \bar{\sigma} \subset H^c$. Lemma 5 [L. 2022] concludes with $X \cap H^c$. □

How to peel fully convex digital sets ?

Context and objectives

Sufficient conditions for peelability

Characterization of peelability

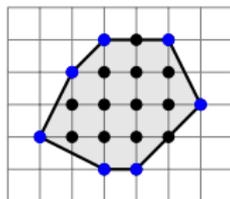
Conclusion and future works

Definition of peelability

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- (1) $z - e_i \in X$ or $z + e_i \in X$, (We say that $i \in \text{Easy}_X[z]$)
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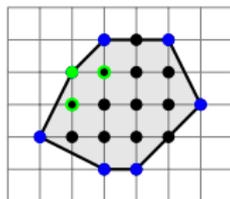


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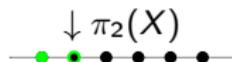
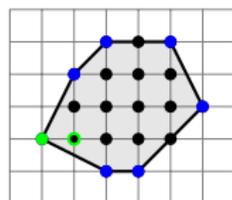


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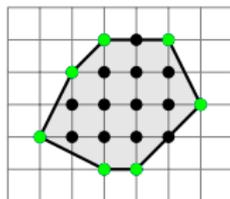


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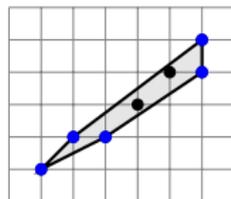
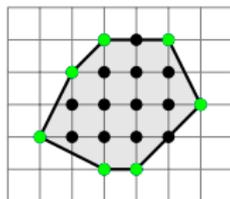


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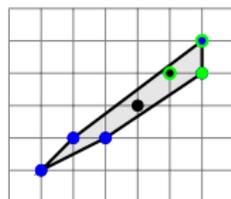
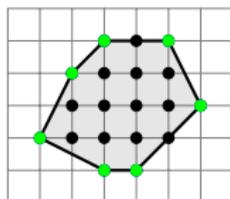


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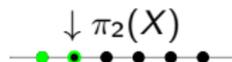
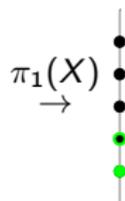
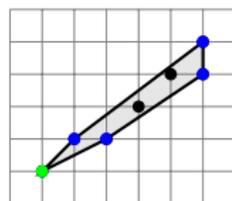
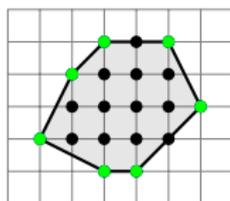


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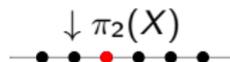
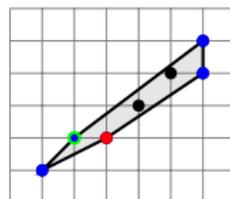
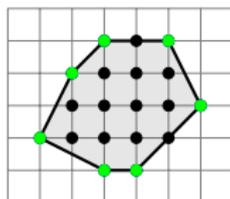


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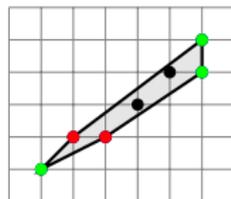
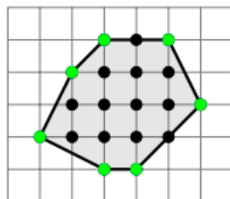


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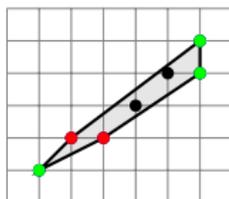
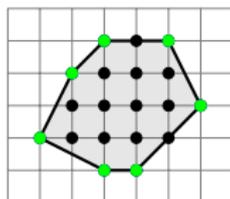


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We postpone here the proof that z is peelable in X iff $X \setminus \{z\}$ is fully convex.

Vertex peelability is locally decidable in 2D

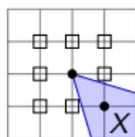
Lemma (vertex peelability is local in 2D)

Let $X \subset \mathbb{Z}^2$, finite and fully convex, and $N_X(z) := \{y \in X, \|z - y\|_\infty \leq 1\}$ the neighborhood of z in X .

A vertex z of $\text{Cvxh}(X)$ is peelable in X iff z is peelable in $N_X(z)$.

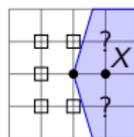
Proof.

By enumeration on $\#(X \cap (z[e_1] \cup z[e_2]))$:



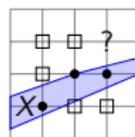
$N = 0$

peelable



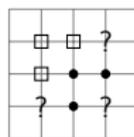
$N = 1$

peelable



$N = 1$

not peelable

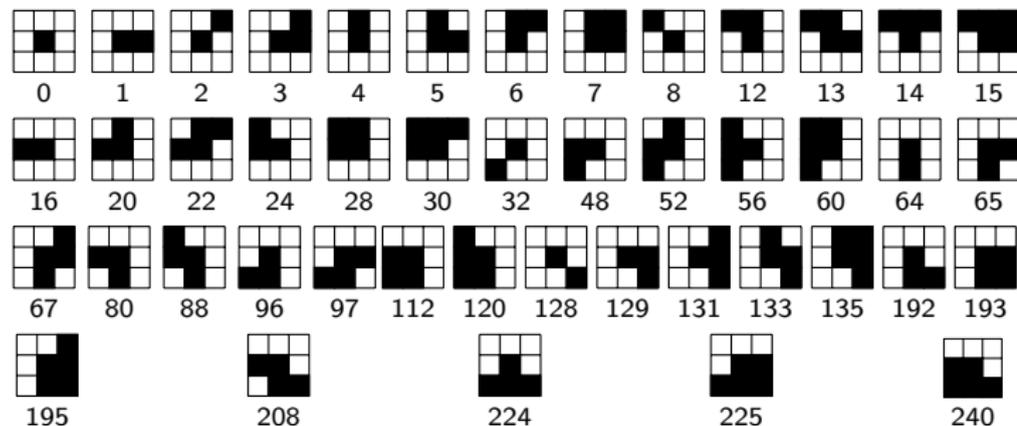


$N = 2$

peelable



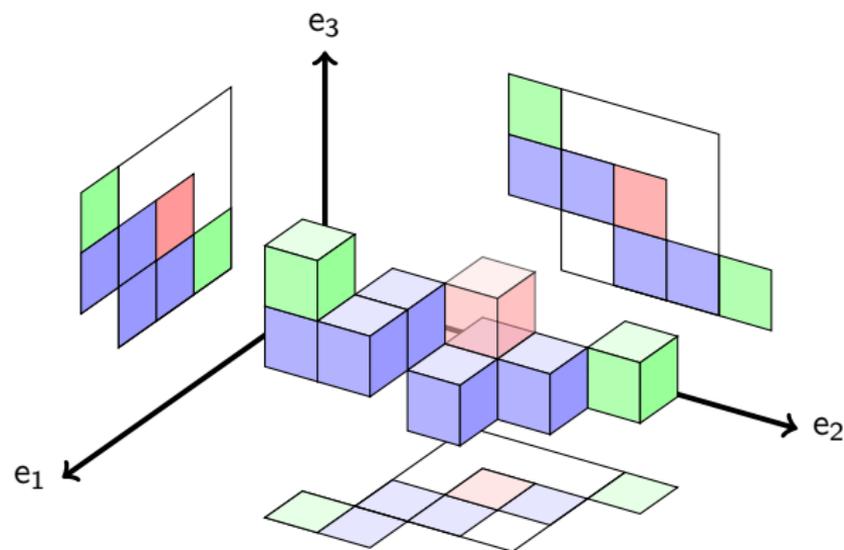
Peelable configurations and algorithm in 2D



3	2	1
4	0	0
5	6	7

```
// input: cfg the encoded neighborhood of the point (in {0,...,255})
// output: true iff the point is peelable wrt its neighborhood
// prerequisite: the point should be a vertex of CvxH(X)
bool isLocallyPeelable( unsigned int cfg )
{ // 45 configurations are peelable
  static const unsigned int p[ 8 ]
    = { 0x5151f1ff, 0x11110001, 0x00101000b, 0x01010003,
        0x000000ab, 0x00000000, 0x00001000b, 0x00010003 };
  return p[ cfg >> 5 ] & ( 1 << (cfg & 0x1f) );
}
```

Peelability is not locally decidable starting from 3D



- ▶ Point z is a vertex of $C_{\text{vxh}}(X)$
- ▶ it is peelable in $\pi_i(X)$ for the 3 directions
- ▶ it is *not* a vertex of $\pi_1(X)$ (not detectable in $N_X(z)$)
- ▶ thus z is **not peelable** in X

Peelability guarantees full convexity

Theorem (peelability guarantees full convexity)

Let X be a non empty finite fully convex set of \mathbb{Z}^d . Then, z is a peelable point of X iff $X \setminus \{z\}$ is fully convex.

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Proof.

\Rightarrow z is peelable in X so z is a vertex of $\text{Cvxh}(X)$. By Lemma 1, $X \setminus \{z\}$ is 0-convex. Then, recursively along projections using P -convexity, if i is easy, $\pi_i(X \setminus \{z\})$ is P -convex. If i is hard, $\pi_i(X \setminus \{z\})$ is 0-convex by peelability and we can go on recursively.



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\Leftarrow X and $X \setminus \{z\}$ fully convex implies z is a vertex of $\text{Cvxh}(X)$. If $\text{Easy}_z[X] = \{1, \dots, d\}$, z is peelable in X . Else, let $i \in \text{Hard}_z[X]$. We look at $\mathcal{Z}_i := z[e_i] \cap \text{Cvxh}(X)$ and **assume** that $\mathcal{Z}_i \neq \{z\}$. Let F^ϵ be farthest facet of $\text{Cvxh}(X)$ along ϵe_i touching \mathcal{Z}_i . It is also a facet of $\text{Cvxh}(X \setminus \{z\})$ and separates z from $z + \epsilon e_i$. The 1-cell $c =]z, z + \epsilon e_i[$ touches F^ϵ and $\text{Cvxh}(X \setminus \{z\})$, so by full convexity $\text{Extr}(c) \cap \text{Cvxh}(X \setminus \{z\}) \neq \emptyset$, a **contradiction**. So $\mathcal{Z}_i = \{z\}$, and this implies that z is a vertex of $\pi_i(X)$.



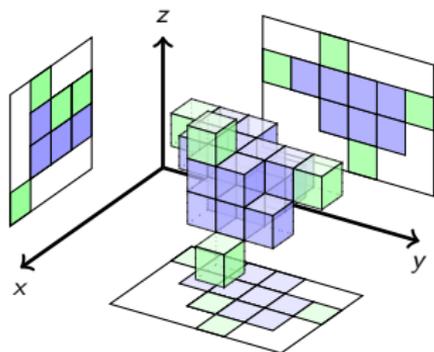
Every fully convex set is peelable

Theorem (every fully convex set is peelable)

Suppose X is a non empty fully convex set in \mathbb{Z}^d . Then, any facet of the bounding box of X contains a peelable point.

Proof (by recursion on the dimension d).

$d = 1$ X has 1 or 2 peelable points.



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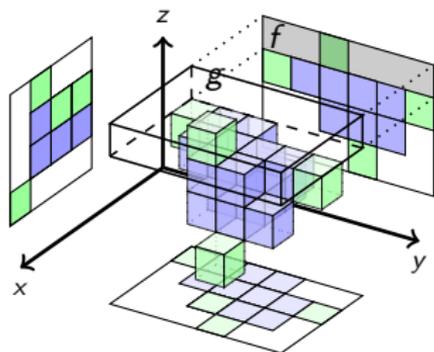
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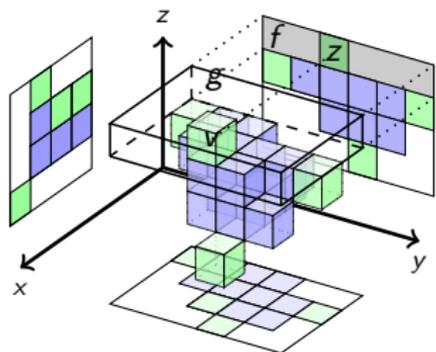
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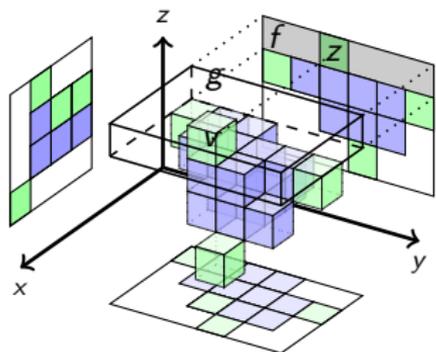
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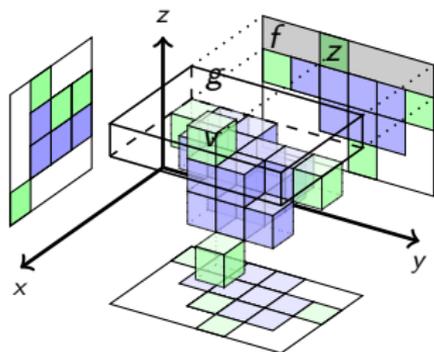
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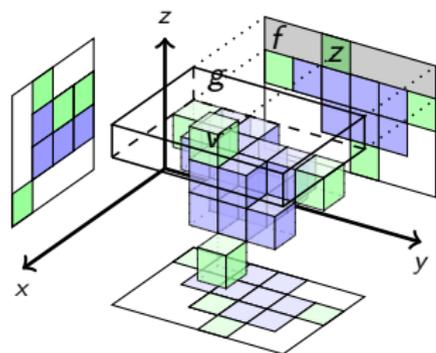
If i easy, we are done. If i is hard, $\pi_i(v) = z$ is a vertex of X_i since it belongs to its bbox .



Every fully convex set is peelable

Theorem (every fully convex set is peelable)

Suppose X is a non empty fully convex set in \mathbb{Z}^d . Then, any facet of the bounding box of X contains a peelable point.



Proof (by recursion on the dimension d).

$d = 1$ X has 1 or 2 peelable points.

$d > 1$ Let f be a facet of $X_i := \pi_i(X)$.

It is a projection of a facet g of $\text{bbox}(X)$.

By recurs. hyp. $\exists z \in f$ peelable point of X_i .

Then $\exists v$ a vertex of $\text{Cvxh}(X)$ with $v \in F$ and $\pi_i(v) = z$.

Hard directions $j \neq i$ for v are stable by projections, but z is peelable, so it remains only i .

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Finally any facet of $\text{bbox}(X)$ is the antecedent of some facet f . □

How to peel fully convex digital sets ?

Context and objectives

Sufficient conditions for peelability

Characterization of peelability

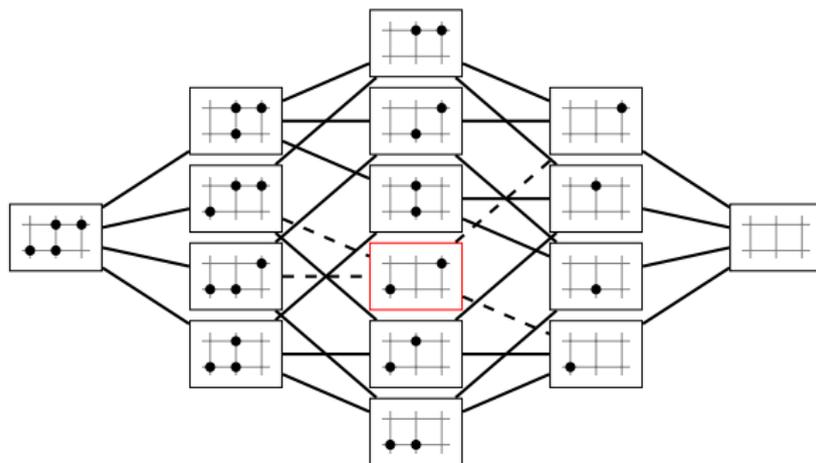
Conclusion and future works

Conclusion

- ▶ better understanding of how to shrink full convex sets
- ▶ peelability is a local property in 2D only
- ▶ recursive characterization of peelability in arbitrary dimension
- ▶ no isolated fully convex set in the hierarchy of fully convex sets

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Future works

- ▶ can we design an envelope operator as thin as $\text{FC}^*(\cdot)$ through peeling, e.g by peeling $\text{Cvxh}(X \oplus [-\frac{1}{2}, \frac{1}{2}]^d) \cap \mathbb{Z}^d$?
- ▶ could it be made increasing ?
- ▶ given $Y \subsetneq X$ fully convex sets, can we peel X keeping Y inside ?
- ▶ how to define the intersection of two fully convex sets ?

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Thank you for your attention !

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Let X be a non empty finite fully convex set of \mathbb{Z}^d . Then, z is a peelable point of X iff $X \setminus \{z\}$ is fully convex.

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Proof.

\Rightarrow z is peelable in X so z is a vertex of $\text{Cvxh}(X)$. By Lemma 1, $X \setminus \{z\}$ is 0-convex. Then, recursively along projections using P -convexity, if i is easy, $\pi_i(X \setminus \{z\})$ is P -convex. If i is hard, $\pi_i(X \setminus \{z\})$ is 0-convex by peelability and we can go on recursively.

□

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- ⇐ X and $X \setminus \{z\}$ fully convex implies z is a vertex of $\text{Cvxh}(X)$.
If $\text{Easy}_z[X] = \{1, \dots, d\}$, z is peelable in X . Else, let $i \in \text{Hard}_z[X]$.
We look at $\mathcal{Z}_i := z[e_i] \cap \text{Cvxh}(X)$ and **assume** that $\mathcal{Z}_i \neq \{z\}$.
Let F^ϵ be farthest facet of $\text{Cvxh}(X)$ along ϵe_i touching \mathcal{Z}_i . It is also a facet of $\text{Cvxh}(X \setminus \{z\})$ and separates z from $z + \epsilon e_i$.
The 1-cell $c =]z, z + \epsilon e_i[$ touches F^ϵ and $\text{Cvxh}(X \setminus \{z\})$, so by full convexity $\text{Extr}(c) \cap \text{Cvxh}(X \setminus \{z\}) \neq \emptyset$, a **contradiction**.
So $\mathcal{Z}_i = \{z\}$, and this implies that z is a vertex of $\pi_i(X)$.

